

Background

vanilladb.org

Why do you need a database system?

To store data, why not just use a file system?

Advantages of a Database System

- It answers queries fast
 - Q1: among a set of blog pages, find those pages written by Ash Ketchum after 2011
 - Q2: among a set of employers, increase the salary by 20% for those who have worked longer then 4 years
- Queries (from multiple users) can execute
 concurrently without affecting each other
- It recovers from crash
 - No corrupt data after restart

Data Model and Queries (1/3)

Q1: among a set of blog pages, find those pages written by Ash Ketchum after 2011

Step1: structure your data by following the *relational data model*

 Identify records (e.g., web pages, authors, etc.) with the same fields in your data and place them into respective tables

blog_pages

blog_id	url	created	author_id	
33981	pokemon.com/	2012/10/31	729	← record
33982	apache.org/	2012/11/15	4412	field

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users		V
user_id	name	balance
729	Ash Ketchum	10,235
730	Picachu	NULL

Data Model and Queries (2/3)

Q1: among a set of blog pages, find those pages written by Ash Ketchum after 2011

```
CREATE TABLE blog_pages (
    blog_id INT NOT NULL AUTO_INCREMENT,
    url VARCHAR(60),
    created DATETIME,
    author_id INT);

INSERT INTO blog_pages (url, created, author_id)
    VALUES ('pokemon.com/...', 2012/09/18, 729);
```

blog_pages

blog_id	url	created	author_id
33981	pokemon.com/	2012/10/31	729
33982	apache.org/	2012/11/15	4412

Data Model and Queries (3/3)

Q1: among a set of blog pages, find those pages written by Ash Ketchum after 2011

Step2: issue queries

```
SELECT b.blog_id
    FROM blog_pages b, users u
    WHERE b.author_id=u.user_id
        AND u.name='Ash Ketchum'
    AND b.created >= 2011/1/1;
```

Advantages of a Database System

- It answers *queries* fast
 - Q1: among a set of web pages, find those pages written by Ash Ketchum after 2011
 - Q2: among a set of employers, increase the salary by 20% for those who have worked longer then 4 years
- Queries (from multiple users) can execute concurrently without affecting each other
- It recovers from crash
 - No corrupt data after restart

Transactions (1/3)

 Each query, by default, is placed in a transaction (tx for short) automatically

```
BEGIN TRANSACTION;
SELECT b.blog_id
    FROM blog_pages b, users u
    WHERE b.author_id=u.user_id
        AND u.name='Ash Ketchum'
        AND b.created >= 2011/1/1;
COMMIT TRANSACTION;
```

Transactions (2/3)

- You can group multiple queries in a transaction optionally
- For example, Steven wants to donate \$100 to Picachu:

```
BEGIN TRANSACTION;

UPDATE users

SET balance=blance-100

WHERE name='Ash Ketchum';

UPDATE users

SET balance=blance+100

WHERE name='Picachu';

COMMIT TRANSACTION;
```

Transactions (3/3)

A database ensures the ACID properties of transactions

Atomicity

 All operations in a transaction either succeed (transaction commits) or fail (transaction rollback) together

Consistency

- After/before each transaction (which commits or rollback), your data do not violate any rule you have set
- E.g., blog_pages.author_id must be a valid users.user_id

Isolation

Multiple transactions can run concurrently, but cannot interfere with each other

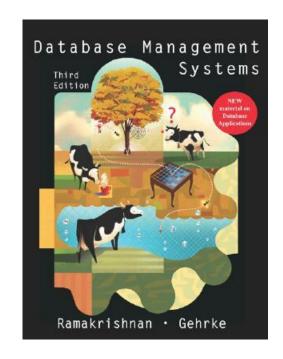
Durability

Once a transaction commits, any change it made lives in DB permanently (unless overridden by other transactions)



Assigned Reading

- Java concurrency
- "Database Management Systems," 3ed, by Ramakrishnan



Coverage

- Java concurrency
- Chaps 2 and 3 on how to store your data into a DBMS
 - ER model and relational model
- Chaps 4 and 5 on queries
 - SQL language (DDL and DML)
 - Relational algebra
- Chap 19* on how to store your data well
 - Easy maintenance
 - Answering most queries fast

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Staring a New Thread

```
public class HelloRunnable implements Runnable {
    public void run() {
        System.out.println("Hello from a thread!");
    public static void main(String args[]) {
        (new Thread(new HelloRunnable())).start();
or
public class HelloThread extends Thread {
    public void run() {
        System.out.println("Hello from a thread!");
    }
    public static void main(String args[]) {
        (new HelloThread()).start();
```

What Happened?

```
public class HelloRunnable implements Runnable {
    public void run() {
        System.out.println("Hello from a thread!");
    }

    public static void main(String args[]) {
        (new Thread(new HelloRunnable())).start();
    }
}
```

- A new stack is allocated in the memory scheme
- Your CPU spends time on executing the code in run ()

Multiple Stacks, Single Heap

- The heap in memory scheme?
 - Stores objects
 - Shared by all threads
- Can two threads access the same object?
 - Yes
- How?
 - Passing the same object to their constructors

Thread Interference

- Given the same object o
- Suppose two threads execute

```
int c = o.get();
c++; // c--;
o.set(c);
```

```
class Counter {
    private int c = 0;
    public void set(int c) {
        this.c = c;
    }
    public int get () {
        return c;
    }
}
```

- Thread A's result will be lost if
 - 1. Thread A: Get c
 - Thread B: Get c
 - Thread A: Increment retrieved value; result is 1
 - Thread B: Decrement retrieved value; result is -1
 - 5. Thread A: Set result in c; c is now 1.
 - Thread B: Set result in c; c is now -1.

Synchronization

```
public class SynchronizedCounter {
  private int c = 0;
  public synchronized void set(int c) {
      this.c = c;
  public synchronized int get() {
      return c;
                     public class SynchronizedCounter {
                        private int c = 0;
                        public void set(int c) {

    Same as

                            synchronized(this) { this.c = c; }
                        public int get() {
                            synchronized(this) { return c; }
```

Memory scheme?

Still Wrong!

 Solution1: the caller locks o during the entire increment/decrement period:

```
synchronized(o) {
   int c = o.get();
   c++; // or c--;
   o.set(c);
}
```

• Solution2: callee provides atomic methods

```
public class SynchronizedCounter {
   private int c = 0;
   public void synchronized increment() {
        c++;
   }
   public synchronized int get() {
        return c;
   }
}
```

Blocking and Waiting

- Threads are blocked outside a critical section if some other is in
- A thread A in a critical section of o can give up the lock by calling o.wait()
 - So, some other blocking thread B can be in
 - A can regain the lock by o.notifyAll() (called by other threads)

```
while (c == 10) { // c has upper bound
    o.wait();
}
C++;
o.set(c);
```

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Note

- DBMS ≠ database
- A database is a collection of your data stored in a computer
- A DBMS (DataBase Management System) is a software that manages databases

Storing Data

- Let's say, you have data in memory to store
- What's the data in memory (heap) look like?
 - Objects
 - References to objects
 - You define classes, the blueprint
- Could we store these objects and references directly?

Data Model

- Definition: A data model is a framework for describing the structure of databases in a DBMS
- Common data models: ER model and relational model
- A DBMS supporting the relational model is called the relational DBMS

Why ER Model?

- Allows thinking your data in OOP way
- Entity
 - An object (or instance of a class)
 - With attributes
- Entity group
 - A class
 - Must define the ID attribute
- Relationship between entities
 - References (has-a relationship)
 - Could be 1-1, 1-many, and many-many

Why Relational Model?

- A realization of ER model
 - Allows queries to be defined and answered
 - Still logic (not how your data stored physically)

Relation

- Realization of 1) an entity group via table; or 2) a relationship
- Fields/attributes as columns
- Records/tuples as rows

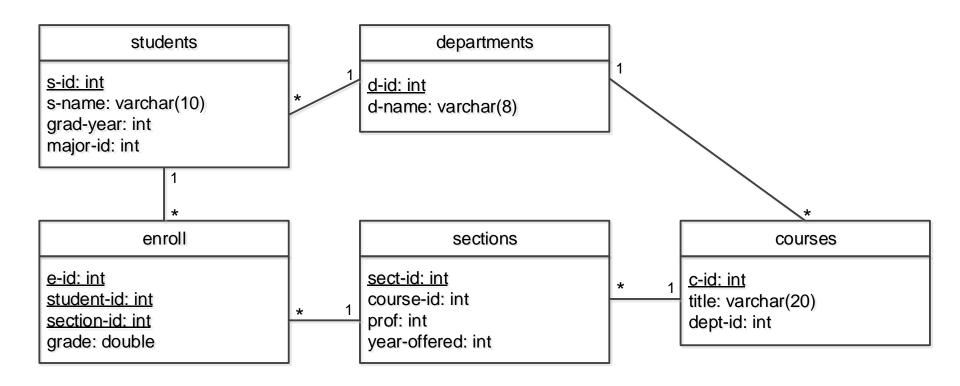
Primary Key

Realization of ID via a group of fields

Foreign key

- Realization of relationship
- A record can have the primary key fields of the other record it refers to
- Only 1-1 and 1-many
- Intermediate relation is needed for many-many

Example: A student DB



Schema

- Definition: A schema is the structure of a particular database
- The schema of a relation/table is its fields and field types

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Queries

- Data Definition Language (DDL) on schema
 - CREATE TABLE …
 - ALTER TABLE …
 - DROP TABLE …
- Data Manipulation Language (DML) on records
 - INSERT INTO ... VALUES ...
 - SELECT ... FROM ... WHERE ...
 - UPDATE ... SET ... WHERE ...
 - DELETE FROM ... WHERE ...

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33982	apache.org/	2012/11/15	4412	field

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users		<u> </u>
user_id	name	balance
729	Ash Ketchum	10,235
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   FROM blog_pages b, users u
WHERE b.author_id=u.user_id
   AND u.name='Ash Ketchum'
AND b.created >= 2011/1/1;
```

How Is a Query Answered?

SELECT b.blog_id

FROM blog_pages b, users u

WHERE b.author_id=u.user_id

AND u.name='Ash Ketchum'

AND b.created \geq 2011/1/1;

product(b, u)

blog_id	url	created	author_id	user_id	name	balance
33981		2009/10/31	729	729	Ash Ketchum	10,235
33981		2009/10/31	729	730	Picachu	NULL
33982		2012/11/15	4412	729	Ash Ketchum	10,235
33982		2012/11/15	4412	730	Picachu	NULL
41770		2012/10/20	729	729	Ash Ketchum	10,235
41770		2012/10/20	729	730	Picachu	NULL

blog id url cre

blog_id	url	created	author_id
33981		2009/10/31	729
33982		2012/11/15	4412
41770		2012/10/20	729

user_id	name	balance
729	Ash Ketchum	10,235
730	Picachu	NULL

How Is a Query Answered?

```
SELECT b.blog_id
```

```
FROM blog_pages b, users u
```

```
WHERE b.author_id=u.user_id

AND u.name='Ash Ketchum'

AND b.created >= 2011/1/1;
```

select(p, where...)

blog_id	url	created	author_id	user_id	name	balance
41770		2012/10/20	729	729	Ash Ketchum	10,235



p = product(b, u)

blog_id	url	created	author_id	user_id	name	balance
33981		2009/10/31	729	729	Ash Ketchum	10,235
33981		2009/10/31	729	730	Picachu	NULL
33982		2012/11/15	4412	729	Ash Ketchum	10,235
33982		2012/11/15	4412	730	Picachu	NULL
41770		2012/10/20	729	729	Ash Ketchum	10,235
41770		2012/10/20	729	730	Picachu	NULL



How Is a Query Answered?

```
SELECT b.blog_id
```

FROM blog pages b, users u

```
WHERE b.author_id=u.user_id

AND u.name='Ash Ketchum'

AND b.created >= 2011/1/1;
```

project(s, select...)

blog_id 41770

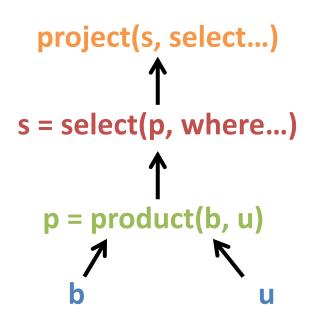
s = select(p, where...)



blog_id	url	created	author_id	user_id	name	balance
41770		2012/10/20	729	729	Ash Ketchum	10,235

Query Algebra

- Operators
 - Product, select, project, join, group-by, etc.
- Operands
 - Tables, output of other operators, predicates, etc.
- Query plan
 - A tree that answers a query
 - Not unique!
- A DBMS automatically seeks for the best query plan



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How Good are Your Data?

- Let's say, if you want to track the topics of a blog page
- Is this a good table?

blog_pages

blog_id	url	created	author_id	topic	topic_admin
33981	pokemon.com/	2012/10/31	729	programming	5638
33981	pokemon.com/	2012/10/31	729	databases	5649
33982	apache.org/	2012/11/15	4412	programming	5638
33982	apache.org/	2012/11/15	4412	os	7423

Insertion Anomaly

blog_pages

blog_id	url	created	author_id	topic	topic_admin
33981	pokemon.com/	2012/10/31	729	programming	5638
33981	pokemon.com/	2012/10/31	729	databases	5649
33982	apache.org/	2012/11/15	4412	programming	5638
33982	apache.org/	2012/11/15	4412	os	7423

33983	apache.org/	2013/02/15	7412	



 A blog cannot be inserted without knowing all fields of topics (except setting them to null)

Update Anomaly

blog_pages

blog_id	url	created	author_id	topic	topic_admin
33981	pokemon.com/	2012/10/31	729	Java prog.	5638
33981	pokemon.com/	2012/10/31	729	databases	5649
33982	apache.org/	2012/11/15	4412	programming	5638
33982	apache.org/	2012/11/15	4412	os	7423

 If you forget to update all duplicated cells, you get inconsistent data

Deletion Anomaly

blog_pages

blog_id	url	created	author_id	topic	topic_admin
33981	pokemon.com/	2012/10/31	729	programmin	5638
33981	pokemon.com/	2012/10/31	729	databases	5649
33982	apache.org/	2012/11/15	4412	programming	5638
33982	apache.org/	2012/11/15	4412	os	7423

 Deleting topics force you to delete the blog fields too

Normalization

- Avoids these anomaly through schema normalization
 - 3rd normal form
 - BCNF normal form
- Idea: break your one, big table into multiple small, modular tables
 - Reuse tables
 - Avoid bias towards any particular pattern of querying